# Two-Way Protocols for occam-π

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#### Before we start...

- This is a proposal
  - It hasn't yet been implemented
- It's a synthesis of several existing ideas
- It's applicable to a variety of process-oriented languages and libraries
  - so when I say "occam", read "occam or JCSP or CHP or PyCSP or ..."

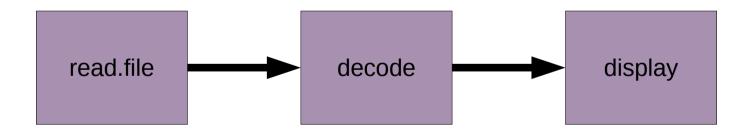


# The problem



#### Processes and channels

 In occam, we build programs by composing processes connected by synchronous, unidirectional channels





#### **Protocols**

- The messages that may be sent over a channel are defined by a protocol
- The compiler checks that the program follows the protocol

```
PROTOCOL POSITION IS INT; INT:

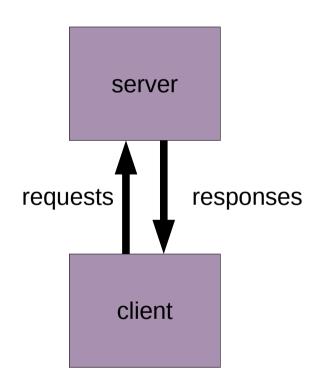
PROTOCOL VIDEO.STREAM

CASE
frame; TIME; [][]PIXEL
end.of.stream
:
```



#### Clients and servers

- A common design pattern: server processes answer requests from client processes
- Design rules can be used to construct complex client-server networks safely





#### Conversations

- Each interaction between a client and server is a conversation, and may contain any number of messages
- For example, the loan pattern:
  - Client: "Let me borrow your big data structure."
  - Server: "OK, here it is."
  - Client: "Right, I'm done; you can have it back now."



#### Client-server in occam

Request and response channels have separate protocols

```
PROTOCOL LOAN.REQ
CASE
borrow
return; MOBILE DATA
:

PROTOCOL LOAN.RESP
CASE
lend; MOBILE DATA
:
```



### Safety assured?

 We can check the protocol on each individual channel

#### But:

- Client: "Let me borrow your big data structure."
- Server: "OK, here it is."
- (Client gets distracted and wanders off.)
- Client: "Let me borrow your big data structure."
- (Boom!)



### What went wrong?

- Each channel's protocol is checked, but the overall conversation is not checked
  - ... so it's possible for the client and server to get into an inconsistent state
- We need a way of describing the two-way protocol that the client and server follow
  - This is useful for documentation too!



# Some existing approaches



#### Honeysuckle (Ian East)

- Language for engineering client-server systems
- A compound service defines the interface to a server using simplified code

```
sequence
receive command
if command
write
acquire String
read
transfer String
```



# Session types (Kohei Honda)

 A formal way of describing two-way communication protocols in terms of the communications that may occur

```
INT! . INT!
(write! . STRING!) | (read! . STRING?)
borrow! . lend? . DATA? . return! . DATA!
```



#### Session types (Honda)

- Originally proposed for use with the pi-calculus
- Several implementations in various languages
  - For concurrency
  - For network protocols



#### State machines

- Session types can be statically checked by translation into finite state machines
- Session type is a (state machine, state ID) pair
- Communications update the state ID



# Proposal



#### Two-way channels

- Add two-way channels to occam-pi
- Can support communication in either direction
  - ... provided both ends agree on the direction
    - You can't ALT between c! and c?
  - Existing channel implementations (CCSP, JCSP et al.)
     already support this
- Superset of existing channel facilities



#### Two-way protocols

- Message content and direction is specified using two-way protocols
  - These are session type declarations
- Conversations must always be started by the same end...
  - so we can always tell what direction the next communication will be in
  - This is already one of the client-server design rules:
     the client must initiate conversation



# Splitting up

 In classical occam, one input/output operation performs the whole one-way protocol

CHAN POSITION c:

c ! 42; 13

**POSITION** protocol



# Splitting up

Now, a two-way protocol may describe several operations

```
CHAN LEND c:
MOBILE DATA thing:
SEQ
```

```
c ! borrow
c ? lend; thing
-- do something with thing
c ! return; thing
```

#### LEND protocol



### Checking the protocol

- The occam compiler can check this by attaching a session type to each channel end
  - ... which is updated on each communication

```
-- c has session type:
-- lend? . DATA? . return! . DATA!
c ? lend; thing
-- c has session type:
-- return! . DATA!
```



# Delegation's what you need

- Since the compiler tracks the session type of each channel end, you can manipulate them safely in the middle of a conversation
  - Abbreviate them
  - Pass them to a procedure
  - For mobile channel ends, communicate them to another process
- Can also split a one-way communication across multiple lines



# Multiple uses

- Can use this to build client-server systems (as in Honeysuckle)
- But it's not tied to the client-server design rules, so it's useful for other types of process network too
- This can replace several existing uses of channel bundles – reduces overhead a bit!



# Syntax



#### Session types in occam

- You'll notice I haven't shown how you define a two-way protocol in occam yet
- There are several possible syntaxes we could consider
- I want to get this right suggestions appreciated!



#### One approach

- Adapt session types notation into occam syntax
  - This is what most session types implementations do
  - Similar to existing one-way protocol syntax

```
PROTOCOL LOAN IS borrow!;

lend?; MOBILE DATA?;

return!; MOBILE DATA!:

PROTOCOL STORE IS (read!; STRING?)

OR (write!; STRING!):
```



#### Another way

- Use simplified occam code
  - ... like Honeysuckle does
  - More verbose, but clearer for complex protocols

```
PROTOCOL LOAN
SEQ
! borrow
? lend; MOBILE DATA
! return; MOBILE DATA
.
```



### The problems

- Both approaches have strengths and weaknesses...
  - Describe the lifetime of the channel, or just a single transaction?
  - Reusing and extending protocols
  - Describing a particular state: LOAN[lend]
  - Elegance and similarity to existing syntax
- See the paper for more details



#### Thanks!

Any questions?

