

Lazy Cellular Automata with Communicating Processes

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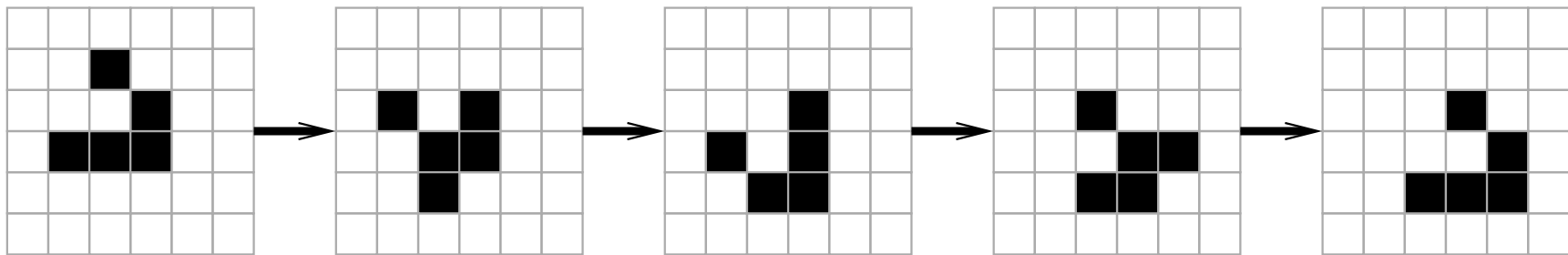
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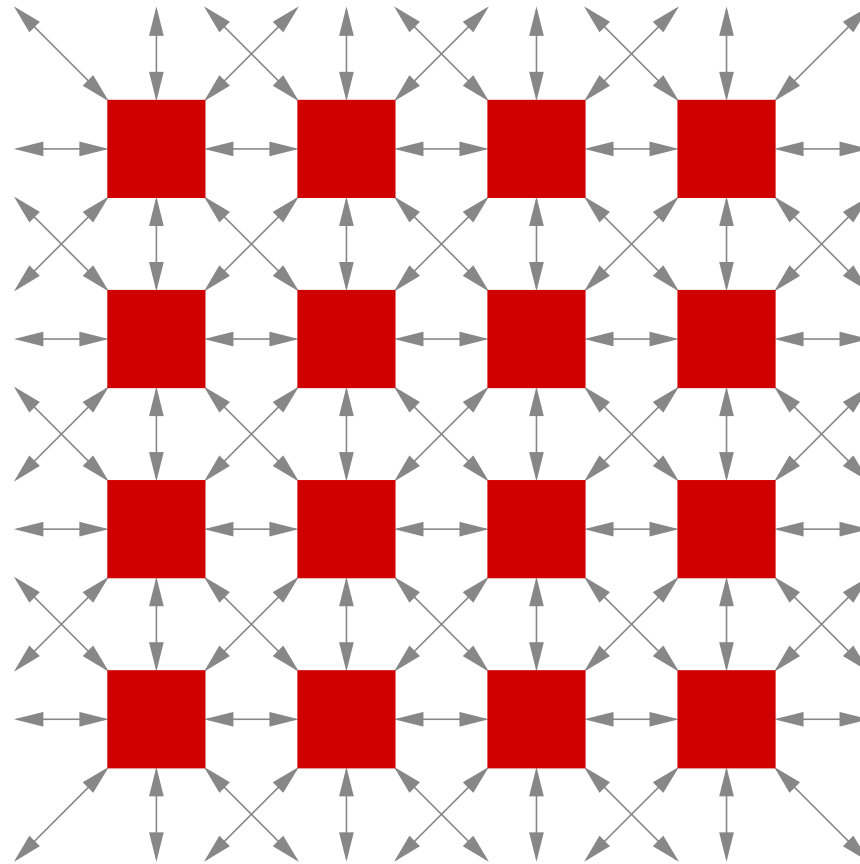
- The Theory Underpinning Nanotech Assemblers project needs to use PC clusters to simulate large numbers of autonomous entities
- ... which we're modelling as CAs for now
- We're using the `occam- π` and JCSP languages – based on CSP extended with some ideas from the π -calculus
- Let's look at Life as an example...

- Infinite grid of cells, each alive or dead
- On each generation step, examine self and 8 adjacent cells
- Alive and 2 or 3 live neighbours \longrightarrow alive
- Dead and exactly 3 live neighbours \longrightarrow alive
- Otherwise \longrightarrow dead
- Interesting emergent behaviour – e.g. the “glider”:



(Black cells are alive.)

- One process per cell, connected in a grid



```
proc cell
  while true
    par i = 0 for 8
      ... send state to neighbour[i]
      ... read state from neighbour[i]
    ... compute new state
```

- Inefficient
- 16 communications per cell per generation
- Most of the time the state hasn't changed
- ... so we only want to communicate changes
- We need a new way of synchronising generation steps

- Barriers synchronise a set of processes
- Processes **sync** on the barrier, and block until *all* the enrolled processes are trying to **sync**...
- ... at which point they all proceed happily
- We can use a barrier for our generation tick

```
proc cell
  ... exchange initial state with
        neighbours (as before)
while true
  ... compute new state
  if my state has changed
    par i = 0 for 8
      ... send state to neighbour[i]
          down buffered channel
sync barrier
  ... check buffered channels
        for changes from neighbours
```

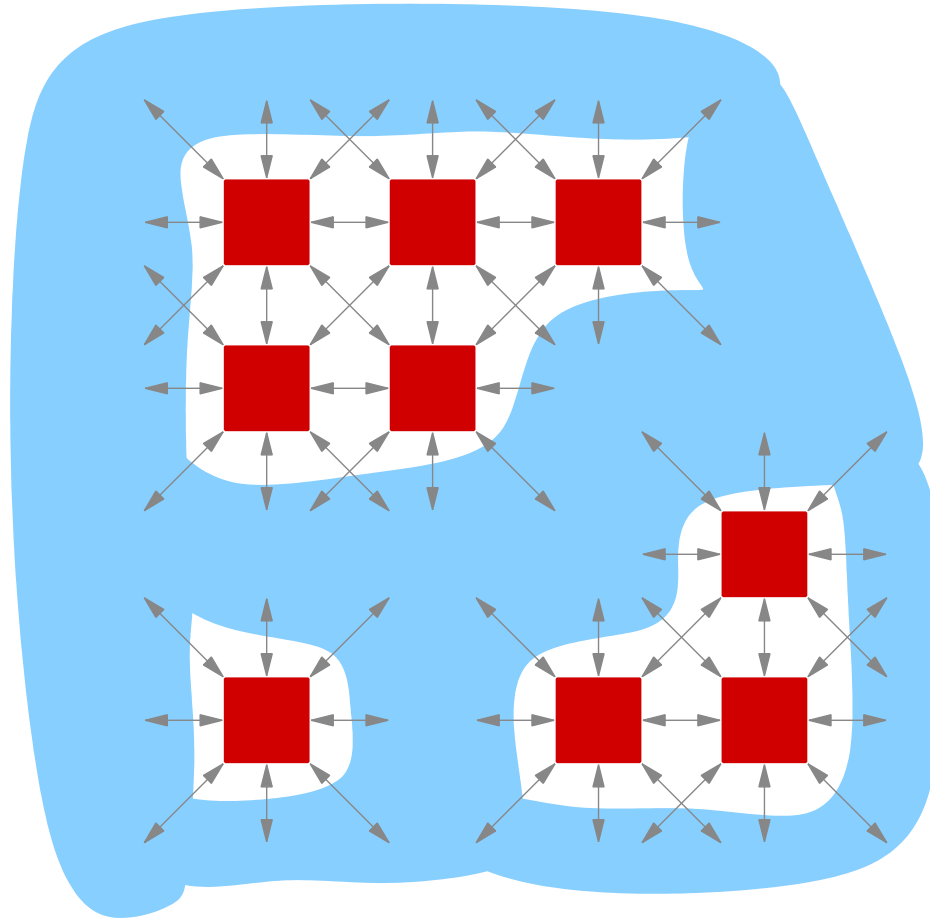

- This is still inefficient
- All cells have to run and synchronise every generation, even if nothing around them has changed
- ... so we want them to “sleep” when possible
- Make them **resign** from the barrier

```
proc cell
  ... exchange initial state
while true
  ... compute new state
  if my state has changed
    par i = 0 for 8
      ... send state to neighbour[i]
    else
      resign barrier
      ... wait for a change to be received
  sync barrier
  ... check for changes
```

- This is *still* inefficient
- Lots of channels – can use one shared channel per cell
- Lots of sleeping processes
- ... so let's only create processes for the active cells
- Use FORKING and a shared state array
- Use phases to control access to the array

```
proc cell
  running := true
  while running
    phase 1 -- state array is constant
    ... read neighbour state from array
    ... compute my new state
    if my state has changed
      ... fork new cell processes for
         the affected neighbours
         (if not already running)
    else
      running := false
  phase 2 -- update state array
  ... write my new state to array
```

- Using shared memory isn't very occam-ish
- Plus we've still got a big array in memory
- ... so use FORKING, but still connect cells with channels
- Use mobile channels to dynamically build the active bits of the network



Clumps of active cells, connected by mobile channels,
floating in the ether

- No particular reason why each process should only simulate one cell
- Make each process simulate a block of cells
- Can take advantage of existing fast sequential code
- ... or mix-and-match parallel approaches

- A number of approaches for simulating CAs in CPA environments
- Same approaches could be applied to other simulation tasks (FEA, CFD)
- Applications for new functionality in *occam- π*
- See the paper for more details – ask us if you'd like a copy of the demo code

Any questions?